

# Least Quadratic Non-Residue assuming Weak GRH

## Weaker Bound

This is a simple analysis to find the Least quadratic non residue in  $\mathbb{F}_p$ . The calculation is much less rigorous and the bound is a bit weaker  $\{O(\log^4(p))\}$  as compared to that of Ankeny [1].

## Notations

The least quadratic non-residue modulo prime  $p$  is denoted by  $n(p)$ .  $\chi_0, \chi_2$  denotes the trivial and quadratic(1 on residues and  $-1$  at non-residues) Dirichlet character respectively.  $\psi(x) = \sum_{i \leq x} \Lambda(i)$  and  $\Lambda(n) = \log p$  if  $n = p^r$  for some prime  $p$  and 0 otherwise.

I will now state few basic facts of Analytic number theory. One can find proof of all these facts in [2].

### Fact 1.

$$\psi(x) = x - \sum_{\rho} \frac{x^{\rho}}{\rho} - \frac{\zeta'(0)}{\zeta(0)} + \frac{1}{2} \log(1 - x^{-2})$$

here  $\rho$  are the non-trivial roots of the Riemann zeta function.

Using Fact 1 it is easy to see that assuming *Riemann Hypothesis*

$$\psi(x) = x + O(\sqrt{x} \log^2 x)$$

**Fact 2.** Let  $\psi(x, \chi) = \sum_{i \leq x} \Lambda(i) \chi(i)$  and  $\chi$  be a non-trivial character mod  $p$

$$\psi(x, \chi) = - \sum_{\rho} \frac{x^{\rho}}{\rho} - \frac{L'(0, \chi)}{L(0, \chi)} + \text{polylog}(p)$$

here  $\rho$  are the non-trivial roots of the Dirichlet L function.

**Theorem 1.** *Assuming GRH,  $n(p) = O(\log^4 p)$ .*

*Proof.* Consider

$$S(M) = \sum_{1 \leq a \leq M} \chi_o(a)\Lambda(a) - \sum_{1 \leq a \leq M} \chi_2(a)\Lambda(a)$$

Note that:  $S(M)$  is positive iff there is a quadratic non residue in  $[1, M]$ .

We know,

$$\begin{aligned} S(M) &= \psi(M, \chi_o) - \psi(M, \chi_2) \\ &= M + O(M^{0.5} \log^2(pM)) \quad (\text{by GRH}) \end{aligned}$$

Since  $S(M) > 0$ ,  $M > CM^{0.5} \log^2(Mp)$ . Therefore  $M = O(\log^4 p)$

□

### Weaker Generalized Riemann Hypothesis

The Generalized Riemann hypothesis says all the non-trivial roots  $\rho$  of the Dirichlet L function are on real part  $z = 0.5$ , but what if we instead of assuming “the exact Riemann hypothesis” consider a weaker form of it. Instead of saying that all the non trivial roots lie on  $RE(\rho) = 0.5$ . We consider that all the non trivial roots lie in a smaller critical strip  $[0.5 - \epsilon, 0.5 + \epsilon]$ . So depending on  $\epsilon$  will our bound will remain poly(log n) or becomes polynomial in n, that is the basic question to be encountered.

Using Fact 2 it is easy to see that

$$\psi(x, \chi) = O(x^{0.5+\epsilon} \log^2(pm))$$

**Theorem 2.** *In Case of weaker Generalized Riemann hypothesis,*

$$n(p) = O(\log^{\frac{4}{1-2\epsilon}} p)$$

*Proof.* Consider

$$S(M) = \sum_{1 \leq a \leq M} \chi_o(a)\Lambda(a) - \sum_{1 \leq a \leq M} \chi_2(a)\Lambda(a)$$

Again  $S(M)$  is positive iff there are no quadratic non residue in  $[1, M]$  We know,

$$\begin{aligned} S(M) &= \psi(M, \chi_o) - \psi(M, \chi_2) \\ &= M + O(M^{0.5+\epsilon} \log^2(pm)) \quad (\text{by weak GRH}) \end{aligned}$$

Since  $S(M) > 0$ ,  $M > CM^{0.5+\epsilon} \log^2(Mp)$ . Therefore  $M = O(\log^{\frac{4}{1-2\epsilon}} p)$

□

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# Bibliography

- [1] Ankeny, N. C. “The Least Quadratic Non Residue.” *Annals of Mathematics Second Series* 55.1 (1952): 65-72. Web.
- [2] Ram Murty, “Problems in Analytic Number Theory, Springer ” , <https://books.google.co.in/>, publisher